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A General Methodology For Equipping Ontologies With Time

Hans-Ulrich Krieger

Deutsches Forschungszentrum für Künstliche Intelligenz GmbH (DFKI)

What Is This Talk All About?

representing changing relationships over time important for

- reasoning & querying services on top of RDF & OWL
- practical applications, e.g., business intelligence
- Semantic Web & Web 2.0 in general
- DLs unable to represent diachronic relations directly
 - no built-in mechanism to handle changing relationships
 - temporal DLs are no exception
 - extending relation instances with time leads to massive proliferation of objects

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- 4D view makes it easy to extend ontologies with time
- preferable: a temporal "annotation" mechanism plus lightweight temporal reasoning services

Example: Synchronic Relation

Tony Blair was born on May 6, 1953.

output of an IE system (RDF triples):

<tb, rdf:type, Person> <tb, hasName, "Tony Blair"> <tb, dateOfBirth, "1953-05-06">

dateOfBirth is a synchronic relation, often functional temporal entity stored as range value of relation instance representation perfectly captures the intended meaning

Example: Diachronic Relation

most relationships vary with time

Christopher Gent was Vodafone's chairman until July 2003. Later, Chris became the chairman of GlaxoSmithKline with effect from 1st January 2005.

informal IE output:

[????-??,2003-07-??]: <cg, isChairman, vf> [2005-01-01,???-???]: <cg, isChairman, gsk>

Example: Diachronic Relation, cont.

applying synchronic representation scheme from above gives:

<cg, isChairman, vf> <cg, hasTime, [????-??, 2003-07-??]> <cg, isChairman, gsk> <cg, hasTime, [2005-01-01, ????-??-??]>

resulting RDF graph mixes up association between fact and extent: [????-??., 2003-07-??]: <cg, isChairman, vf> [2005-01-01, ????-??-??]: <cg, isChairman, vf> [????-??-??, 2003-07-??]: <cg, isChairman, gsk> [2005-01-01, ????-??-??]: <cg, isChairman, gsk>

Encoding 1: Equip Relation with Temporal Argument

obvious extension, used in temporal data bases and logic programming community

$$hasCeo(c, p) \longmapsto hasCeo(c, p, \underline{t}) \text{ or } hasCeo(c, p, \underline{s, t})$$

DLs do **not** support relations with more than two arguments, i.e., encoding **not** applicable to OWL

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Side Note: Temporal Description Logics

so what are Temporal Description Logics (e.g., Lutz 2004)?

TDLs = DLs + concrete domain (Baader & Hanschke 1991)TDLs are great aiming at representing **synchronic** relations temporal features are **functional** relations descriptive inventory: paths, additional constructors (e.g., <)

example:

Human $\sqsubseteq \exists$ (hasMother.dateOfBirth < dateOfBirth) $\sqcap \exists$ (hasFather.dateOfBirth < dateOfBirth)

Encoding 2: Apply Meta-Logical Predicate

use **holds** to encode temporally constant information hasCeo must be reinterpreted as a functional fluent used by situation calculus, Allen logic, KIF complex relation arguments not possible in OWL annotation properties in OWL not possible for relation instances

 $hasCeo(c, p, t) \longmapsto holds(hasCeo(c, p), t)$

Encoding 3: Reify Original Relation

relation reification loses original relation needs introduction of a new class for each relation requires massive ontology rewriting new individual, four additional relation instances similarities to reification in RDF

 $\begin{array}{l} \underline{hasCeo}(c, \ p, \ t) \ \longmapsto \ \exists hc \, . \\ type(hc, \ \underline{HasCeo}) \land hasTime(hc, \ t) \land \\ company(hc, \ c) \land person(hc, \ p) \end{array}$

Encoding 4: Wrap Range Arguments

domain argument often anchor for reasoning and querying so wrap range arguments in a new container object same container class can be applied to each relation instance ontology rewriting still needed

related to relation reification, but does not lose relation name

 $\begin{array}{ll} \mathsf{hasCeo}(\mathsf{c}, \ \mathsf{p}, \ \mathsf{t}) &\longmapsto \ \exists \mathsf{et} \ .\\ \mathsf{type}(\mathsf{et}, \ \overline{\mathsf{EntityTime}}) \land \mathsf{hasTime}(\mathsf{et}, \ \mathsf{t}) \land \\ \mathsf{hasCeo}(\mathsf{c}, \ \mathsf{et}) \land \mathsf{hasEntity}(\mathsf{et}, \ \mathsf{p}) \end{array}$

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Perdurants and Time Slices: Encoding 5+6

distinction between *endurants* and *perdurants* in philosophy perdurantist view: all entities only exist for some period of time perdurant \approx 4D trajectory in spacetime time slice = temporal part of a 4D slice of special interest: slices where specific information stays constant we usually only have partial information for a given perdurant

Encoding 5: Encode Perdurantist/4D View in OWL

Welty & Fikes 2006: OWL implementation of perdurantist view time slice encodes time dimension of spacetime relations from source ontology no longer connect original entities encoding requires ontology rewriting

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\begin{array}{l} \underline{\mathsf{hasCeo}(\mathsf{c},\,\mathsf{p},\,\mathsf{t})} \longmapsto \exists \mathsf{ts1},\,\mathsf{ts2}\,.\\ \\ \overline{\mathsf{hasCeo}(\mathsf{ts1},\,\mathsf{ts2})} \land \\ \\ \mathrm{type}(\mathsf{ts1},\,\mathsf{TimeSlice}) \land \mathsf{hasTimeSlice}(\mathsf{c},\,\mathsf{ts1}) \land \mathsf{hasTime}(\mathsf{ts1},\,\mathsf{t}) \land \\ \\ \\ \mathrm{type}(\mathsf{ts2},\,\mathsf{TimeSlice}) \land \mathsf{hasTimeSlice}(\mathsf{p},\,\mathsf{ts2}) \land \mathsf{hasTime}(\mathsf{ts2},\,\mathsf{t}) \end{array}
```

Encoding 6: Reinterpret Perdurantist/4D View

reinterpret perdurantist view:

what has originally been an entity becomes a time slice original entities now describe the "behavior" of perdurants at a certain moment in time (e.g., being a person)

time slices of a perdurant need not to be of the same type, e.g., perdurant DFKI has slices for Company and AcademicInstitution cooccurring information in such a slice stays constant encoding does NOT need rewriting of original ontology

 $\begin{array}{ll} \mathsf{hasCeo}(\mathsf{c}, \mathsf{p}, \underline{\mathsf{t}}) &\longmapsto \\ \mathsf{hasCeo}(\mathsf{c}, \mathsf{p}) \land \mathsf{hasTime}(\mathsf{c}, \mathsf{t}) \land \mathsf{hasTime}(\mathsf{p}, \mathsf{t}) \land \\ \mathsf{hasTimeSlice}(\mathsf{C}, \mathsf{c}) \land \mathsf{hasTimeSlice}(\mathsf{P}, \mathsf{p}) \end{array}$

time slices c, p are linked to perdurants C, P (created only once)

Example I

DC's CEO Jürgen Schrempp announces that he will resign by 31st December 2005.



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Example II

I believe [that] Jürgen Schrempp was the CEO of DC from 1995 until 2005.



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Equiping OWL Ontologies With Time: Example

- 1. find out which relations will undergo a temporal change
- 2. identify domain and range class(es) for these relations
- 3. make these classes time slices using owl:equivalentClass

example: PROTON upper ontology (proton.semanticweb.org/)

- 1. most properties in PROTON are diachronic properties
- 2. psys:Entity is the class of choice, both for domain and range
- 3. fourd:TimeSlice \equiv psys:Entity



General Integration Scheme

1. always use 4D

Perdurant: hasTimeSlice; TimeSlice: timeSliceOf, hasTime

2. choose Time

an arbitrary time ontology, e.g., OWL-Time

3. **choose upper/domain ontology** the original ontology that lacks time, e.g., PROTON

4. use Allen (optional)

13 relations, plus 6 super-relations defined over time slices

add axiom fourd:TimeSlice ≡ c₁ □ ... □ c_n
 c₁,..., c_n: maximal incompatible classes that need to be extended by a temporal dimension

Outlook: Temporal Extensions to OWL

additional arguments, going beyond binary relations/triples Hayes-/ter Horst-style rules can be extended by a temp. dimension only lightweight reasoning needed

example 1: owl:inverseOf ceoOf(js, dc, 1995, 2005) → hasCeo(dc, js, 1995, 2005)

example 2: owl:SymmetricProperty
marriedWith(bbt, aj, 2000, 2003)
→ marriedWith(aj, bbt, 2000, 2003)

example 3: owl:TransitiveProperty contains(dfki, room1.26, s, t) & contains(room1.26, chair42, u, v) \rightarrow contains(dfki, chair42, *max*(s, u), *min*(t, v))

Paper: Further Issues

- sophisticated time ontology
 - temporal underspecification
 - granularity of time
- more on Hayes-/ter Horst-style entailment rules
- comparison how extended tuples ease the writing of custom rules (and querying), compared to RDF triples

Thank you!

Questions?

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